



De-anonymization attack on geolocated data

Sébastien Gambs, Marc-Olivier Killijian, Miguel Nuñez del Prado Cortez

► To cite this version:

Sébastien Gambs, Marc-Olivier Killijian, Miguel Nuñez del Prado Cortez. De-anonymization attack on geolocated data. *Journal of Computer and System Sciences*, 2014, 80 (8), pp.1597-1614. 10.1016/j.jcss.2014.04.024 . hal-01242268

HAL Id: hal-01242268

<https://hal.science/hal-01242268>

Submitted on 11 Dec 2015

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

De-anonymization attack on geolocated data

Sébastien Gambs^{a,*}, Marc-Olivier Killijian^b, Miguel Núñez del Prado Cortez^{b,c}

^aUniversité de Rennes 1 - INRIA / IRISA, Campus Universitaire de Beaulieu 35042 Rennes, France

^bLAAS - CNRS, 7 avenue du Colonel Roche, BP 54200, F-31031 Toulouse, France

^cUniversité de Toulouse, INSA, LAAS, F-31400 Toulouse, France

Abstract

With the advent of GPS-equipped devices, a massive amount of location data is being collected, raising the issue of the privacy risks incurred by the individuals whose movements are recorded. In this work, we focus on a specific inference attack called the de-anonymization attack, by which an adversary tries to infer the identity of a particular individual behind a set of mobility traces. More specifically, we propose an implementation of this attack based on a mobility model called Mobility Markov Chain (MMC). A MMC is built out from the mobility traces observed during the training phase and is used to perform the attack during the testing phase. We design several distance metrics quantifying the closeness between two MMCs and combine these distances to build de-anonymizers that can re-identify users in an anonymized geolocated dataset. Experiments conducted on real datasets demonstrate that the attack is both accurate and resilient to sanitization mechanisms such as downsampling.

Keywords: Privacy, geolocation, inference attack, de-anonymization.

1. Introduction

With the recent advent of ubiquitous devices and smartphones equipped with positioning capacities such as GPS (Global Positioning System), a massive amount of mobility traces is collected and gathered in the form of geolocated datasets by cellphone companies, providers of location-based services, and developers of smartphone applications. Some of these geolocated datasets are available in public repositories and can be used for research or for industrial purposes (*e.g.*, to optimize the placement of cellular towers, to conduct market and sociological studies or to analyze the flow of traffic inside a city). These datasets are composed of the mobility traces of hundreds or thousands of individuals [1, 2, 3, 4], thus raising the issue of the privacy risks incurred by these individuals. For instance, from the movements of an individual it is possible to infer his points of interests (such as his home and place of work) [5, 6, 7, 8], to predict his past, current and future locations [9, 10], or even to infer his social network [11].

In this work, we focus on a particular form of inference attack called the *de-anonymization attack*, by which an adversary tries to infer the identity of a particular individual behind a set of mobility traces. More precisely, we suppose that the adversary has been able to observe the movements of some individuals during a non-negligible amount of time (*e.g.*, several days or weeks) in the past during the *training phase*. Later, the adversary accesses a different geolocated dataset containing the mobility traces of some of the individuals observed during the training phase, plus possibly some unknown persons. Then, the objective of the adversary is to de-anonymize this dataset (called the *testing dataset*) by linking it to the corresponding individuals contained in the training dataset. Note that simply replacing the real names of individuals by pseudonyms before releasing a dataset is usually not sufficient to preserve the anonymity of their identities because the mobility traces themselves contain information that can be uniquely linked back to an individual. In addition, while a dataset can be sanitized before being released by adding spatial and temporal noise, the risk of re-identification through de-anonymization attack nevertheless still exists. Thus, in order to be able to assess the importance of this risk, it is important to develop a method to quantify it.

*Corresponding author

Email addresses: sgambs@irisa.fr (Sébastien Gambs), marco.killijian@laas.fr (Marc-Olivier Killijian), mnunezde@laas.fr (Miguel Núñez del Prado Cortez)

In this paper, we propose a novel method to de-anonymize location data based on a mobility model called Mobility Markov Chain (MMC) [7]. A MMC is a probabilistic automaton, in which each state corresponds to one (or possibly several) Point Of Interest (POI), characterizing the mobility of an individual and an edge indicates a probabilistic transition between two states (*i.e.*, POIs). Each state can have an semantic label attached to it such as “home”, “work”, “leisure”, “sport”, ... A MMC is built out of the mobility traces observed during the training phase and is used to perform the de-anonymization attack during the testing phase. More precisely, the mobility of each individual both from the training and testing sets is represented in the form of a MMC. Afterwards, a distance is computed between possible pairs of MMCs from the training and testing sets in order to identify the closest individuals in terms of mobility. In short, the gist of this method is that the mobility of an individual can act as a signature, thus playing the role of a quasi-identifier [12]. Thus, if the adversary knows a target user Alice and a signature of her mobility (*e.g.*, he has learnt her MMC out of the training set), he can try to identify her by finding a matching signature in the testing set.

The outline of the paper is the following. First, in Section 2, we review some related work on de-anonymization attacks and mobility models before briefly introducing in Section 3 the background on Mobility Markov Chains necessary to understand our work. Afterwards in Section 4, we present the distance metrics between MMCs that we designed in order to quantify the closeness between two mobility behaviors, while in Section 5 we describe how to build predictors (which we call *de-anonymizers*) based on these distances to efficiently and accurately de-anonymize location data. Finally, we evaluate experimentally the efficiency of the proposed attack on real geolocated datasets in Section 6 before concluding in Section 7.

2. Related work

An *inference attack* corresponds to a process by which an adversary that has access to some data related to individuals (and potentially some auxiliary information) tries to deduce new personal information that was not explicitly present in the original data. For instance, a famous inference attack was conducted by Narayanan and Shmatikov on the “Netflix dataset” [13]. This dataset is a sparse high dimensional data containing ratings on movies of more than 500 000 subscribers from Netflix that was supposed to have been anonymized before its release for the Netflix competition¹. However, Narayanan and Shmatikov have performed a de-anonymization attack that was able to successfully re-identify more than 80% of the Netflix subscribers by using the Internet Movie Database² (IMDB), a public database of movie ratings, as auxiliary knowledge.

Inference attacks have also been developed specifically for the geolocated context. For instance, Mulder *et al.* [6] have proposed two methods for profiling users in a GSM network that can also be used to perform a de-anonymization attack. The first method is based on constructing a Markov model of the mobility behavior of an individual while the second considered only the sequence of cell IDs visited. Once POIs have been extracted for each user, an agglomerative hierarchical clustering algorithm is used to group users according to a similarity measure called the *cosine similarity* [14]. Their first method is relatively similar to ours with two major differences due to the fact that their dataset comes from a cellular network. Thus, it relies on the static GSM cells as the states of the Markov model, while we dynamically learn the POIs from the mobility traces of an individual. Therefore, in our setting two individuals do not necessarily have any POI in common, whereas with GSM cells, individuals living in the same area have a high probability of sharing some POIs (*i.e.*, cells). As a consequence, the second main difference between their work and ours is that the transitions are only possible between neighboring GSM cells, as it is impossible to “jump” from one cell to another if they are not adjacent. Their attack was validated through experimentations using cell locations of 100 users from a MIT Media Lab dataset³ [15] that includes information such as call logs, Bluetooth devices in proximity, cell tower IDs, application usage and phone status. During the experiments, the authors have observed that if the currently profiled user belongs to a cluster of other similar users, there is a high chance of making a mistake about the identity inferred among all the users of this cluster when performing the de-anonymization attack. The success rate of the re-identification attack varies from 37% to 39% using the Markov model, against 77% to 88% when the sequence of cells visited is used.

¹<http://www.netflixprize.com>

²<http://www.imdb.com>

³<http://reality.media.mit.edu/dataset.php>

Zang and Bolot [8] have performed a study of the top n most frequently visited places by an individual in a GSM network and show how they can act as quasi-identifiers to re-identify anonymous users. Their study was performed on the Call Data Records (CDRs) from a nationwide US cellular provider collected over a month and contains approximately 20 millions users. From this dataset, the authors have identified the top n most frequent places for each user at different levels of spatial (*i.e.*, sector, cell, zip code, city, state and country) and time granularity (*i.e.*, day and month). Their inference attack was able to successfully re-identify 35% of the population studied when the adversary has no auxiliary knowledge and even up to 50% when the adversary can use the knowledge of the social network of users as auxiliary information. The social network was constructed by creating a social relationship between two individuals that have called each other at least once in the past. In their analysis, the authors emphasize that the distance between home and work can be an indicator of the privacy level for an individual. In particular, the larger this distance, the higher the risk that this individual can be de-anonymized.

Ma *et al.* [16] proposed an inference attack to de-anonymize users in a geolocated dataset along with a metric to quantify the privacy loss of an individual. Two datasets were used in this study, one recording the movements of San Francisco YellowCabs [1] and the other related to the movements of Shanghai city public buses⁴. Two types of adversary models were considered: the passive one, collecting the whereabouts of individuals from a public source (possibly sanitized) and the active one that can deliberately participate or perturb the data collection in order to gain additional knowledge about the location of some specific individuals. To retrieve the identities of individuals, the authors imagine four different estimators that the adversary can use to measure the similarity between mobility traces (for instance between the original traces and the sanitized ones). Namely, the Maximum Likelihood Estimator relies on the Euclidean distance to compute the similarity between mobility traces. The Minimum Square Approach computes the negative sum of the square of the absolute value of the difference between the original traces and the sanitized ones. The Basic Approach assumes that the traces have been perturbed by uniform noise. Thus with this approach, a mobility trace will be considered to be similar to another trace if it is contained within a sphere of radius r centered on the original trace. Finally, the Weighted Exponential Approach is similar to the Basic Approach, with the exception that no assumption is made on the type of noise generated. These methods approach a success rate of de-anonymization of 80% to 90% on the San Francisco YellowCabs dataset and between 60% and 70% on the Shanghai bus dataset, and this even when the data is sanitized through the addition of spatial noise. However, contrary to our work, these inference attacks were conducted on the whole dataset (there was no split between a training set and a testing set). In particular, the authors assume that both type of adversaries (*i.e.*, passive and active ones) pick the information they need to build the mobility model from the same dataset on which the success of the de-anonymization attack is tested. This induces an overly strong bias in the re-identification results obtained with this approach.

The work of Freudiger, Shokri and Hubaux [17] also focuses on re-identifying users of geolocated datasets. These experiments have been conducted on two datasets. The first dataset⁵ contains the GPS traces of 24 users from the city of Borlange recorded in a two-year period (1999-2001) while the second one is due to Nokia⁶ and is composed of the GPS traces of 150 users from the city of Lausanne recorded over a year. In this work, the pair of POIs “home/work” is used as pseudo-identifiers to de-anonymize users. First, a variant of the k -means algorithm is used to extract POIs from the mobility traces. Then, the POI in which the individual considered stays the most often between 9PM to 9AM is identified as “home” while the POI in which the individual stays the most often between 9AM to 5PM is labelled as “work”. The training phase consists in applying this method on the raw data to extract the pairs “home/work” for all individuals, and then to conduct the same attack on sampled traces in order to assess how much the pair “home/work” can still be inferred even when the dataset released has been sanitized by applying sampling. The success of this method depends on how sampled traces have been generated. Indeed, the authors have proposed several sampling schemes whose bias towards selecting home/work locations or other POIs can be parametrized. The authors have shown that when 100 samples are observed, the de-anonymization rate is approximately 67% for the Borlange dataset and 70% for the Nokia one. Unfortunately, as most of the previous works presented, this study does not split the available data into a training and testing set during its evaluation but rather generates the samples that will constitute the testing set directly from the training one. This introduces a major bias in the evaluation of the techniques, which we further discuss in Section 6.4.

⁴To the best of our knowledge, this dataset is not publicly available contrary to the other one.

⁵<http://icapeople.epfl.ch/freudiger/borlange.zip>

⁶To the best of your knowledge, this dataset is not publicly available.

The work of Xiao *et al.* [18] relies on the notion of Semantic Location Histories (*SLH*) to compute the similarity between users. In a nutshell, a *SLH* is simply the sequence of semantic locations frequently visited by an individual. Like several previous approaches, this work first uses a hierarchical clustering algorithm to extract POIs out of mobility traces. Then, using as external knowledge a database that can associate semantics to a location, each POI is linked with a semantic tag for each level of the hierarchy (e.g., “Italian restaurant” and then “restaurant” at an upper level). Finally, the *SLH* is computed by analyzing the sequence of POIs visited by a user and taking into account their semantic labels. The similarity measure designed by the authors is based on the notion of *maximal travel match*, which counts the number of similar semantic locations visited (not necessarily in the same order) by two different *SLH*s within a predefined time interval. This metric is computed for each layer of the cluster hierarchy before being summed over all possible layers, possibly by weighting the influence of a particular level (*i.e.*, the deeper the level, the bigger the influence). Finally, the proposed approach was evaluated on the Geolife dataset [19]. Contrary to previous work, the success of the de-anonymization attack is quantified in terms of the *normalized discounted cumulative gain*, a metric originated from information retrieval [20]. In a nutshell, the objective of this metric is to rank all the possible candidates to de-anonymization with respect to how close their mobility is to the behavior of the user considered. In the experiments conducted, the success of the attack as measured by this metric was between 0.7 and 0.9. Basically the closer this value is to 1, the more effective the attack is. Note that, due to the different metric that was used, this method is not directly comparable to other previous works.

Shokri *et al.* [21] inferred the correspondence between pseudonymized traces of 40 randomly chosen users of the YellowCabs dataset [1], in which each position is a cell of 8×5 grid over the San Francisco Bay area, and user profiles represented in the form of hidden Markov model. Their attack computes a matching probability between pseudonymized traces and user profiles by using the classical forward-backward algorithm [22]. This matching probability basically represents the likelihood that a particular set of traces correspond to a specific user. Once these probabilities are computed, a bipartite graph of traces and users is formed whose edges are weighted according to the associated matching probability between pairs of user/traces. This bipartite graph is then given as input to the Hungarian algorithm that assign a particular set of traces to each user. As the paper mainly focuses on solving the task of “individual tracking”, which corresponds to being able to predict at which location a user was located at a particular moment in time, their results are not directly comparable to our work. However, we believe that the framework they have developed is generic enough in terms of the attacks it encompasses that it could also model the de-anonymization attack.

The study of Srivatsa and Hicks [23] de-anonymizes users based on the contact graph, generated from mobility traces and information about their social network. The authors use three different datasets: the Saint Andrews dataset, which contains 18 241 contact traces from 27 different users issued from WiFi access points, the Smallblue dataset is formed by 240 665 contact traces of 125 users representing messages generated by instant messaging, and finally the Infocom 06 dataset that consists of 182 951 contact traces from 78 users extracted from Bluetooth devices used during the Infocom 2006 conference. In addition, the authors have used some auxiliary knowledge in the form of social networks extracted from Facebook for the Saint Andrews and Smallblue datasets and from DBLP for the Infocom 06 dataset. Their attack relies on the structural similarities between contact and social graph, which is measured in terms of standard graph similarity measures such as the graph edit distance (*i.e.*, minimal number of edges to erase or add to have a perfect match), the maximum common subgraph (*i.e.*, the number of vertices in the largest common subgraph) and the distribution of node degrees. The de-anonymization attack works in two steps. During the first step, the n most important POIs are extracted from both contact and social graphs based on the node centrality metric [24]. The second step is the de-anonymization process itself that links the identities of users from the social graph with the corresponding users of the contact graph. The authors have used three different mapping techniques: distance vector (82% of de-anonymization rate), spanning trees (88% of de-anonymization rate) and subgraph matching (80% of de-anonymization rate).

Sharad and Danezis [25] have de-anonymized the communication graph datasets of the D4D challenge⁷. Their attack exploits the topology of the contact graph to break the anonymity by looking for subgraphs that are isomorphic and form the largest subgraphs. More precisely, the authors compute the 1-hop node neighborhood degree distribution for each node of the subgraphs generated from the graph. This distribution is used as a similarity metric to identify the

⁷<http://www.d4d.orange.com>

closest node with the same characteristics in the other graph. Afterwards, for matching the 2-hop nodes, the authors rely on two metrics: the neighbor match, which is the number of common 1-hop nodes they are associated to, and the signature match, which is a metric computed from the in and out degree distributions. To evaluate their methodology, Sharad and Danezis have used the EU mail communication dataset [26] and artificially generated a dataset with similar characteristics to the communication graphs of the D4D challenge obtaining between 56% to 96% of re-identification rate.

A recent study of De Montjoye and co-authors [27] has shown that if an individual has a unique pattern in the anonymized dataset then this is enough to characterize him even if a particular dataset does not contain personal information such as name, age and address. The objective of this study was precisely to measure the uniqueness of 1.5 millions users in an anonymized mobile phone dataset collected from April 2006 to June 2007. By taking 2 and 4 random locations to characterize the mobility patterns of individuals, the authors were able to identify respectively 50% and 95% of the users in the anonymized dataset. However, this study has some limitations as for instance it is not clear how the knowledge that is learnt can be used to later in the future de-anonymize an individual. More precisely, the study shows that during a particular period (*e.g.*, a day or a week) if we consider the sequence of a constant number locations visited by an individual compared to the rest of the population, this sequence is likely to be unique. Nonetheless, this do not preclude the possibility that other individuals might visit the same sequence of locations during another period (*e.g.*, the next day or the next week). Thus, it is not obvious how to build an efficient de-anonymizer from the results of this study. Moreover as for other previous works, there is no separation being made between the training and the testing sets in this study.

Name	Location data	Strategy	Re-identification rate
Location profiling GSM [6]	Call data records	Markov chain Sequence of visited places	37%-39% 77%-88%
Most frequented POIs [8]	Call data records and social network	Top n visited locations	35%-50%
Similarity between trajectories [16]	GPS traces	Similarity of trajectories	60%-90%
Pair home/work [17]	GPS traces	Pair home/work used as quasi-identifiers	67%-70%
Semantic location histories [18]	GPS traces	Semantic trajectories	70%-90% (NDCG)
Likelihood between traces and users [21]	GPS traces	Hungarian algorithm	Not applicable
Linking nodes [23]	GPS traces and social network	Structural similarities between contact and social graphs	82%-88%
Structure of the graph [25]	Communication graph	Structural similarities between graphs	56%-96%
Random locations [27]	Call data records	Use of random locations as quasi-identifiers	50%-95%
Summary statistics [28]	Connections to access points	Similarity between spatial distributions of users	43%-71%

Table 1: Summary of related work.

Finally, a recent work by Unnikrishnan and Naini [28] shows that revealing statistics about the movements of users (instead of directly their mobility traces) can also be used to de-anonymize efficiently geolocated datasets and thus may also be harmful to privacy. To perform the de-anonymization attack, the authors computed through the Hungarian algorithm the minimum weight matching [29] in a bipartite graph, composed of an anonymous and labeled graph, in which nodes represent users and the weights on edges correspond to the distance in term of entropy between the spatial and temporal distribution of nodes across different places. The dataset used is a two weeks logs obtained from the access points from the EPFL campus, using the first week for training and the second one for testing. Using

this technique, the authors obtained 70.5% of success for the de-anonymization.

In this section, we have reviewed the previous work on de-anonymization attacks in the geolocated context, which we summarize in Table 1. Note that the results from some of these studies cannot be directly compared due to the difference in nature of the datasets used, the way the success metrics are defined for an attack and also the fact that in most of these studies the data has not been split in proper training and testing sets. We discussed further this fundamental issue in Section 6.4. In the following sections, we will present our approach to de-anonymization. More precisely, we first introduce how to model the mobility of an individual in the form of a MMC, before describing how to measure the similarity between two MMCs using the distances we propose. Finally, we demonstrate experimentally how to use these distance metrics to perform a de-anonymization attack.

3. Mobility Markov Chain

A *Mobility Markov Chain* (MMC) [7] models the mobility behavior of an individual as a discrete stochastic process in which the probability of moving to a state (*i.e.*, POI) depends only on the previously visited state and the probability distribution on the transitions between states. More precisely, a MMC is composed of:

- A set of states $P = \{p_1, \dots, p_n\}$, in which each state is a frequent POI (ranked by decreasing order of importance), with the exception of the last state p_n that corresponds to the set composed of the union of all infrequent POIs. POIs are learned by running a clustering algorithm on the mobility traces of an individual. These states are associated to a location, and generally they also have an intrinsic semantic meaning. Therefore, semantic labels such as “home” or “work” can often be inferred and attached to them.
- *Transitions*, such as $t_{i,j}$, represent the probability of moving from state p_i to state p_j . A transition from one state to itself is possible if the individual has a non-null probability from moving from one state to an occasional location before coming back to this state. For instance, an individual can leave home to go to the pharmacy and then come back to his home. In this example, it is likely that the pharmacy will not be extracted as a POI by the clustering algorithm, unless the individual visits this place on a regular basis and stays there for a significant amount of time.

Note that several mobility models based on Markov chains have been proposed in the past [7, 30], including the use of hidden Markov models for extracting the semantics of POIs [31]. In a nutshell, building a MMC is a two-steps process. During the first phase, a clustering algorithm is run to extract the POIs from the mobility traces. For instance in the work of Gambs *et al.* [7], a clustering algorithm called Density-Joinable Cluster (*DJ-Cluster*) was used (we rely on the same algorithm in this work), but of course other clustering algorithms are possible. In the second phase, the transitions between those POIs are computed and incorporated in the MMC.

DJ-Cluster takes as input a trail of mobility traces and four parameters: the minimal number of points *MinPts* needed to create a cluster, the maximum radius r of the circle within which the points of a cluster should be contained, the minimal number of days that the traces falling within a cluster should cover and d a threshold distance used to decide if two neighboring clusters should be merged. DJ Cluster works in three phases. During the first phase, which corresponds to a preprocessing step, all the mobility traces in which the individual is moving (*i.e.*, whose speed is above a small predefined value) as well as subsequent static redundant traces are removed. As a result, only static traces are kept. The second phase consists in the clustering itself: all remaining traces are processed in order to extract clusters that have at least *MinPts* points within a radius r of the centre of the cluster. Finally, the last phase merges all clusters that have at least a trace in common or whose medoids are within d distance of each other. Once the POIs (*i.e.*, the states of the Markov chain) are discovered, the probabilities of the transitions between states can be computed. To realize this, the trail of mobility traces is examined in chronological order and each mobility trace is tagged with a label that is either the number identifying a particular state of the MMC or the value “unknown”. Finally, when all the mobility traces have been labeled, the transitions between states are counted and normalized by the total number of transitions in order to obtain the probabilities of each transition. A MMC can either be represented as a transition matrix or as a graph in which nodes correspond to states and arrows represent the transitions between along with their associated probability. When the MMC is represented as a transition matrix of size $n \times n$, the rows and columns correspond to states of the MMC while the value of each cell is the probability of the associated transition between the corresponding states.

4. Distances between mobility Markov chains

In this section, we propose four different distances quantifying the similarity between two mobility Markov chains. These distances are based on different characteristics of the MMCs and thus give different but complementary results. We will rely on these distances in the following sections to perform the de-anonymization attack.

4.1. Stationary distance

The intuition behind the *stationary distance* is that the distance between two MMCs can be evaluated as to the sum of the distances between the closest states of both MMCs. In order to compute the stationary distance, the states of the MMCs are paired in order to minimize this distance. As a result, it is possible that a state from the first MMC is paired with several states of the second MMC (this is especially true if the MMCs are of different size). Furthermore, the computation of the stationary distance heavily relies on the stationary vectors of the MMCs. In a nutshell, the stationary vector of a MMC is a column vector V , obtained by multiplying repeatedly a vector initialized uniformly V_{init} by the MMC transition matrix M until convergence (*i.e.*, until the distribution of values in this vector reaches the stationary distribution of the MMC).

The stationary distance is directly computed from the stationary vectors of two MMCs (hence its name). More precisely, given two MMCs, M_1 and M_2 , the stationary vectors, respectively V_1 and V_2 , of each model are computed. Afterwards, Algorithm 1 is run on these two stationary vectors. For each state in V_1 , the algorithm searches for the closest state in V_2 (lines 5 to 11, Algorithm 1) and then multiplies the distance between these two states by the corresponding probability of the stationary vector of the state of V_1 currently considered (line 12, Algorithm 1).

Once the algorithm has taken into account all states from V_1 , the current value computed represents the distance from M_1 to M_2 ($distance_{AB}$ in line 1, Algorithm 2). This distance is not symmetric as such and therefore in order to symmetrize it, Algorithm 1 is called once again, but on V_2 and V_1 in order to obtain the distance from M_2 to M_1 ($distance_{BA}$ of line 2, Algorithm 2). The result is made symmetrical by computing the average of these two distances (line 3, Algorithm 2).

Algorithm 1 Stationary_distance(V_1, V_2)

```

1:  $distance = 0$ 
2: for  $i = 1$  to  $n_1$  (the number of nodes in  $V_1$ ) do
3:    $MinDistance = 100000$  kilometers
4:   Let  $p_i$  be the  $i^{th}$  node of  $V_1$ 
5:   for  $j = 1$  to  $n_2$  (the number of nodes in  $V_2$ ) do
6:     Let  $p_j$  be the  $j^{th}$  node of  $V_2$ 
7:      $CurrentDistance = \text{Euclidean\_Distance}(p_i, p_j)$ 
8:     if ( $CurrentDistance < MinDistance$ ) then
9:        $MinDistance = CurrentDistance$ 
10:    end if
11:  end for
12:   $distance = distance + Prob_{V_1}(p_i) \times MinDistance$ 
13: end for
14: return  $distance$ 
```

Algorithm 2 Symmetric_stationary_distance(V_1, V_2)

```

1:  $distance_{AB} = \text{Stationary\_distance}(V_1, V_2)$ 
2:  $distance_{BA} = \text{Stationary\_distance}(V_2, V_1)$ 
3:  $distance = (distance_{AB} + distance_{BA})/2$ 
4: return  $distance$ 
```

4.2. Proximity distance

The intuition behind the *proximity distance* is that two MMCs can be considered close if they share “important” states. For instance, if two individuals share both their home and place of work they should be considered as being highly similar. Moreover, the importance of a state is directly proportional to the frequency at which it is visited. Therefore, the first states ordered by decreasing order of importance are compared, then the second ones, then the third ones, and so on. The proximity score obtained for sharing the first states is considered twice as important as the score for sharing the second states, which is itself twice as important as the sharing of the third states, and so forth.

Given two MMC models M_1 and M_2 (ordered in a decreasing manner with respect to their stationary probabilities), this distance is parametrized by a threshold Δ and a *rank*. The objective of the rank is to quantify the importance of matching two states at a specific level. In particular, the higher is the value of the *rank*, the bigger is the weight that will be given to these POIs. For instance for the first pair of POIs, we have set *rank* = 10.

Algorithm 3 starts by verifying for each pair of nodes between M_1 and M_2 if the Euclidean distance between them is less than the threshold Δ (line 8, Algorithm 3). If this condition is met, the value of *rank* is added to the score value (line 9, Algorithm 3). Afterwards, *rank* is divided by two (lines 11-14, Algorithm 3). Once all the pair of nodes have been processed, the global distance is set to be the inverse of the global score if this score is non-null (lines 17-19, Algorithm 3). Otherwise, the distance outputted is set to a large value (e.g., 100 000 kilometers).

Algorithm 3 Proximity_distance($V_1, V_2, \Delta, rank$)

```

1: Sort the states of  $V_1$  by decreasing order of frequency
2: Sort the states of  $V_2$  by decreasing order of frequency
3: score = 0
4: for  $i = 1$  to  $\min(n_1, n_2)$  do
5:   Let  $p_a$  be the  $i^{th}$  node of  $V_1$ 
6:   Let  $p_b$  be the  $i^{th}$  node of  $V_2$ 
7:   distance = Euclidean_distance( $p_a, p_b$ )
8:   if (distance <  $\Delta$ ) then
9:     score = score + rank
10:   end if
11:   rank = rank/2
12:   if (rank = 0) then
13:     rank = 1
14:   end if
15: end for
16: distance = 100000
17: if (score > 0) then
18:   distance = 1/score
19: end if
20: return distance

```

4.3. Matching distance

The *matching distance* is similar to the stationary distance in the sense that this distance also corresponds to the sum of distances between the states of these two MMCs. However, the pairing of the states between the MMCs is done in a different way as each state from the first MMC is paired with one and only one state of the second MMC.

The matching distance is based on the Hungarian method [32], which is a polynomial-time combinatorial optimization algorithm for assignment problems. More precisely, the method works as follows. Given two MMCs, M_1 and M_2 and their corresponding stationary vectors V_1 and V_2 , Algorithm 4 first verifies if the number of nodes in both models is the same. If not, we assume without loss of generality that M_2 is the MMC with the fewer number of states. Before continuing the computation of the distance, “dummy” states are added to M_2 . Each dummy state is a copy of the centroid of the states of M_2 (line 4, Algorithm 4). Once this process is completed, the number of states in V_1 and V_2 is the same.

Algorithm 4 Matching_distance(V_1, V_2)

```
1: Let  $D$  be the distance matrix  $D$  and  $Min$  the minimization matrix
2: Let  $n_1$  be the number of nodes in  $V_1$ 
3: Let  $n_2$  be the number of nodes in  $V_2$  (we suppose that  $V_2$  is the stationary vector with the fewest number of states
   if the number of states between  $V_1$  and  $V_2$  is different)
4: Add fake states to  $V_2$  if necessary to ensure that  $n_2 = n_1$ 
5: for  $i = 1$  to  $n_1$  do
6:   for  $j = 1$  to  $n_2$  do
7:      $D_{i,j} = \text{Euclidean\_distance}(p_i, p_j)$ 
8:     if ( $p_i$  or  $p_j$  is fake) then
9:        $Min_{i,j} = 100000$ 
10:    else
11:       $Min_{i,j} = D_{i,j}$ 
12:    end if
13:  end for
14: end for
15:  $Index = \text{HungarianAssignment}(Min)$ 
16:  $Dist = 0$ 
17: for  $i = 0$  to  $n_1$  do
18:    $row = Index_{i,0}$ 
19:    $col = Index_{i,1}$ 
20:   if ( $V_2.\text{node}(col)$  is not fake) then
21:      $proba = (V_1(row) + V_2(col))/2$ 
22:      $Dist = Dist + D_{row,col} \times proba$ 
23:   else
24:      $proba = (V_1.\text{stationaryVector}(row) + 0)/2$ 
25:      $Dist = Dist + \text{NearestState}(V_1.\text{node}(row), V_2.\text{nodes})$ 
26:   end if
27: end for
28: return  $Dist$ 
```

Afterwards, the *distance matrix* ($D_{i,j}$) and the *minimization matrix* ($Min_{i,j}$) are computed. The distance matrix contains the Euclidean distance between each pair of nodes in M_1 and M_2 (lines 5-14, Algorithm 4), in which nodes in M_2 form the rows of the matrix and nodes in M_1 correspond to the columns of $D_{i,j}$. For instance, the distance $D_{i,j}$ is the Euclidean distance between the node i in M_2 and the node j in M_1 . The minimization matrix is equivalent to the distance matrix, except that each distance $Min_{i,j}$ in which node i is a dummy is assigned a large default value, such as 100 000 kilometers (line 9, Algorithm 4). In short, the objective of this large value is to guide the Hungarian method towards giving priority to real states instead of dummy ones.

The minimization matrix, which is squared, is passed as input to the Hungarian method. The Hungarian method then computes the optimal assignment between states from M_1 and M_2 that minimizes the global sum of distances between pair of states and such that each state from V_1 is exactly matched with one state of V_2 (and *vice-versa*). This optimal matching is returned as a matrix *Index* composed of two columns. Each column contains respectively the indexes of rows and columns of the optimal assignment for the minimization matrix $Min_{i,j}$. Figure 1 illustrates graphically an optimal assignment between two MMCs (one for Alice and one for Bob). The matrix *Index* corresponding to Figure 1 is $I = \{\{1, 1\}, \{2, 3\}, \{3, 2\}, \{4, 3\}, \{5, 4\}\}$.

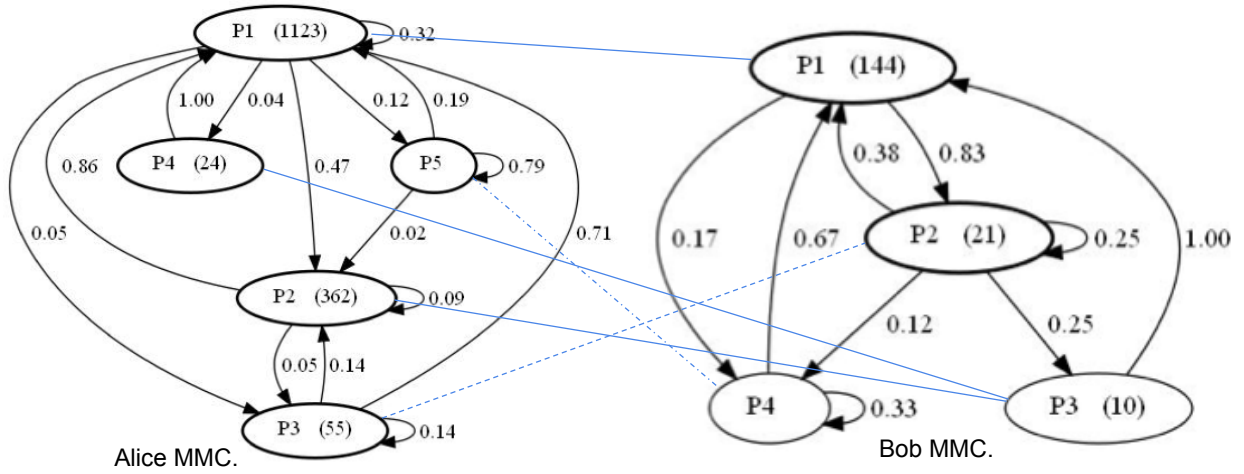


Figure 1: Example of the optimal matching of two MMCs. The number in each node corresponds to the numbers of traces that falls inside this state.

For each pair of the matching obtained, the distance between these two (non-dummy) states is multiplied by the average of probabilities of stationary vectors V_1 and V_2 . However, when one of the state in M_2 is “dummy”, the nearest real state in M_2 is identified, and the distance between these two states is then multiplied by the probability corresponding to the “orphan” state in V_1 divided by two (lines 17–27, Algorithm 4). The distance returned is the sum of the values computed for each pair of the matching (line 28, Algorithm 4).

4.4. Density-based distance

Similarly to the stationary and the matching distances, the *density-based distance* is simply the sum of the distances between pairs of MMCs states. However, the MMCs states are paired according to their rank once they are sorted using their corresponding probabilities in the stationary vector.

First, given two MMC models M_1 and M_2 , the nodes in both models are sorted by decreasing density (lines 1 and 2, Algorithm 5). Therefore, the first node will be the one with the highest stationary probability. Afterwards, the first node of M_1 is matched with the first node of M_2 and the same goes for the rest of the nodes. Finally, the sum of all distances between matched nodes is computed (line 4 to 8, Algorithm 5) and the algorithm outputs the total distance.

The stationary distance, the matching distance and the density-based distance are all numerical, in the sense that they are formed of the sum of the Euclidian distance between some pairing of states. In contrast, the proximity distance is different as it is based on the semantics behind the MMCs. Indeed, the first state in the model is inferred as being very representative of the mobility of an individual (*e.g.*, home), the second as quite representative (*e.g.*, the place of

Algorithm 5 Density – based_distance(V_1, V_2)

```
1: Sort the states of  $V_1$  by decreasing order of density
2: Sort the states of  $V_2$  by decreasing order of density
3:  $distance = 0$ 
4: for  $i = 1$  to  $\min(n_1, n_2)$  do
5:   Let  $p_a$  be the  $i^{th}$  node of  $V_1$ 
6:   Let  $p_b$  be the  $i^{th}$  node of  $V_2$ 
7:    $distance = distance + \text{Euclidean\_distance}(p_a, p_b)$ 
8: end for
9: return  $distance$ 
```

work), and two individuals are considered as very similar if they share these two places. Table 2 summarizes the main characteristics of these distances. Thereafter, we will see how these distance can be used to build de-anonymizers and how their diversity can be exploited and combined to enhance the success of the de-anonymization attack.

Name	Strategy	Type	Complexity
Stationary distance	Stationary vector	Numerical	$O(n)$
Proximity distance	Proximity of MMCs states	Relative	$O(n)$
Matching distance	Hungarian algorithm	Numerical	$O(n^3)$
Density-based distance	Weight of MMC states	Numerical	$O(n)$

Table 2: Summary of MMC distances. In this table, the (computational) complexity is measured with respect to n the number of individuals in the datasets that has to be de-anonymized (for simplicity the training and testing datasets are assumed to contained the same number of individuals).

5. De-anonymizers

In this section, we rely on the four distances proposed in the previous section to build statistical predictors in order to perform a de-anonymization attack. We call such a predictor, a *de-anonymizer* in reference to its main objective. A de-anonymizer takes as input the MMC representing the mobility of an individual and tries to identifies within a set of anonymous MMCs, the one that is the most similar (*i.e.*, the closest in terms of distance). For example, a de-anonymizer may learn from the training set a MMC representing the mobility of Alice and later look for the presence of Alice in the testing set. A de-anonymizer can be based on one distance or a combination of them.

- *The minimal distance* de-anonymizer considers that in each row (*i.e.*, a row refers to an individual to de-anonymize), the MMC with the minimal distance (*i.e.*, the column) is the individual corresponding to the identity of the row. We have considered four instantiations of this de-anonymizer, one for each of the distance previously described (*i.e.*, stationary distance, proximity distance, matching distance and density-based distance).
- *The maximal-gap* de-anonymizer takes as input several distance matrices corresponding to different distance metrics. For each distance matrix, the minimal distance de-anonymizer outputs two predictions (instead of one) for each row, which corresponds to the first and second smallest values of the distances in this row. Afterwards, the gap (*i.e.*, difference) between these two values are computed. The higher the gap, the more confident we can be in the prediction made by the de-anonymizer on this particular distance matrix. Therefore, the distance metric with the highest gap among all the ones considered will be the candidate considered for the de-anonymization.
- *The simple vote* de-anonymizer receives as input a list of candidates, which corresponds to the identities of the n minimal values of a particular row in at least two different distance matrices. The candidate that receives the highest number of votes will be the one considered for the de-anonymization. For example, in Table 3, each column corresponds to different distance metrics while each row contains the proposed candidates at this position. In the first row, Bob gets two votes against one for Alice and therefore he is considered as the most likely candidate for de-anonymization.

Points	m_1	m_2	m_3
3	Bob	Alice	Bob
2	Charlie	Bob	Alice
1	Alice	Charlie	Charlie

Table 3: Illustration of the voting method, each column corresponds to the output of a de-anonymizer based on a different distance metric.

- *The weighted vote* de-anonymizer, much like the simple vote one, takes as input a list of candidates coming from different distance matrices. However, the voting method weights each possible candidate depending on the rank he has obtained for the different distance metrics. For instance, if each distance metric proposes n candidates, the first one is given a weight n , while the second receives a weight of $n - 1$, and so on. For example, in Table 3, the first column contains the weights for each candidate. In this particular case, the outcome of the weighted voting method is that Bob gets 8 votes, Alice 6 and Charlie 4. Therefore, Bob is considered as the “winner” (*i.e.*, most likely candidate) for the de-anonymization.
- *The stat-prox* de-anonymizer behaves exactly like the minimal stationary distance de-anonymizer, except when the stationary distance is above a given threshold and the proximity distance is below its maximum value (*i.e.*, 100 000 kilometers). The intuition is that if the minimal stationary distance is very small, we should use it. Otherwise, we rely on the minimal proximity distance unless it gives no conclusive result, in which case we roll back to the minimal stationary distance.

6. Experiments

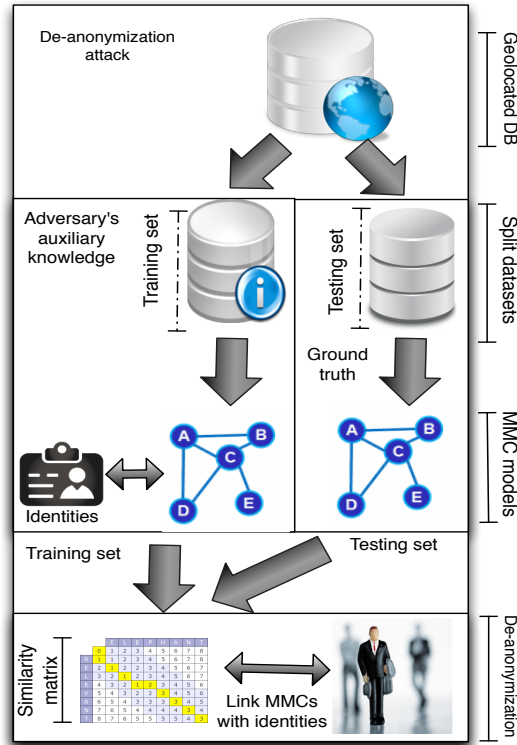


Figure 2: Overview the de-anonymization attack process.

An overview of the process of *de-anonymization attack* over geolocated datasets used in our experiments is illustrated in Figure 2. Considering a particular geolocated dataset, we first sort the mobility traces of each user in a chronological order. Then, for each user his trail of mobility traces is split into two disjoint trail of traces of same size, one for the training set and one for the testing set as explained in Section 6.1. The former is part of the auxiliary knowledge gathered by the adversary while the latter is the ground truth we use to assess the success of the attack. For each of this dataset, we learn a MMC for each trail of mobility traces. With respect to the MMCs learnt from the training set, the adversary knows the correspondence between these models and the corresponding identities of the users. Afterwards, the distances described in Section 4 are used to compute a distance matrix between the MMC models resulting from the training and testing sets. Subsequently, using this distance matrix as input to one of the de-anonymizer described in Section 5, the objective of the de-anonymization attack is to map back the users of the testing set to their true identities by linking their models in the testing set to the corresponding ones in the training set. Finally, the success rate of the attack is computed by measuring the ratio between the number of correct predictions over the total number of guesses. In order to estimate the robustness of the de-anonymization attack when the sampling conditions change, we have also performed some downsampling on the data. In a nutshell, a sampling mechanism summarizes several mobility traces into fewer traces. This is generally done by representing a set of subsequent traces that have occurred within the same time window into a single trace (*e.g.*, the average or median trace).

We evaluate the efficiency of the de-anonymization attack introduced in the previous section on five different datasets described in Section 6.1. Then, we describe the method used to fine-tune the parameters of the clustering algorithms in Section 6.2. Afterwards, we report the results of those experiments that were conducted by using the distances and de-anonymizers described in the previous sections. More precisely, we evaluate the accuracy of the de-anonymization attack relying on either a single predictor or a combination of them (Section 6.3). Finally, in Section 6.4, we compare our work with the performance reported in related works, highlighting in particular the bias in some of the experimental settings of these previous works.

6.1. Description of datasets

The datasets used in the experiments are described thereafter.

1. The *Arum dataset* [4] is composed of the GPS traces of 5 researchers sampled at a rate of 1 to 5 minutes in the city of Toulouse from October 2009 to January 2011.
2. The *Geolife dataset* [19] has been gathered by researchers from Microsoft Asia and consists of GPS traces collected from April 2007 to October 2011, mostly in the area of Shanghai city. This dataset contains the mobility traces of 178 users captured at a very high rate of 1 to 5 seconds.
3. The *Nokia dataset* [2] is the result of a data collection campaign performed the city of Lausanne for 200 users started in September 2009 and that lasted for more than two years. The rate at which the location has been sampled varies depending on the current battery level.
4. The *San Francisco Cabs (SFC) dataset* [1] contains GPS traces of approximately 500 taxi drivers collected over 30 days, between May and July 2008, in the San Francisco area.
5. The *Borlange dataset* [17] has been collected as a part of traffic congestion experiment over two years from 1999 to 2001. The public version of this dataset contains the GPS traces of 24 vehicles.

Table 4 summarizes the main characteristics of the datasets described above, namely the total number of users in the dataset, the collection period measured in the number of days as well as the average number of traces per user and the total number of traces in the dataset.

Characteristics	Arum	Geolife	Nokia	SFC	Borlange
Total nb of users	6	175	185	165	24
Collection period (nb of days)	255	146	1443	23	4103
Average nb of traces/user	16374	121 615	63145	20812	12555
Total nb of traces	98249	21 203 955	11 681 896	34 341 00	301 332

Table 4: Main characteristics of the datasets used.

Conducting an analysis of the distribution of the average number of traces per user, we have observed that there is a high variance with a large gap between the users with a small number of traces compares to the one with a high number of traces. For instance in Figure 3, we have shown the distribution of the users according to their number of traces for the Geolife and Nokia datasets. While these distributions are peaked around the category of users that have between 1000 and 10000 traces, the other categories of users also contains a significant number of users. With respect to the other datasets (Arum, SFC and Borlange), they are close to be uniformly distributed with the exception of the SFC dataset, which also displays a peak distribution.

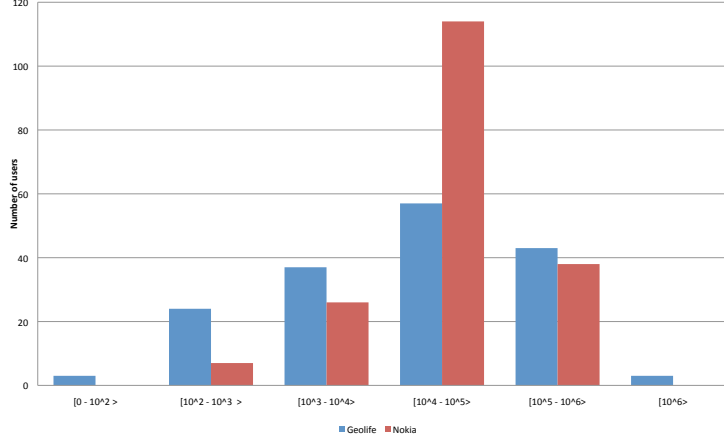


Figure 3: Distribution of the users according to their number of traces for the Geolife and Nokia datasets.

For each individual, we split his trail of mobility traces (chronologically ordered) into two disjoint parts of approximately the same size. The first half of the original data forms the training set, and will be used as the adversary background knowledge, while the second half constitutes the testing set on which the de-anonymization attack is conducted. For instance, if the original trail of one individual is composed of n mobility traces $\{mt_1, mt_2, \dots, mt_n\}$, it will be split into a training set $\{mt_1, mt_2, \dots, mt_{\frac{n}{2}}\}$ and a testing set $\{mt_{\frac{n}{2}+1}, mt_{\frac{n}{2}+2}, \dots, mt_n\}$ (for illustration purpose we assume that n is an even number). Therefore, the objective of the adversary is to de-anonymize the individuals of the testing set by linking them to their corresponding counterparts in the training set. For simplicity reason, we assume thereafter that the training and the testing set are composed of the same n persons while in general this might not be the case (*e.g.*, the testing set could only contain a small fraction of the users of the training set). Of course, dividing the training and test sets based on the number of traces may not guarantee that the length (*i.e.*, number of days) of the period they covered is exactly the same.

In order to evaluate the impact of the size of the training set on the construction of a MMC, we have varied its size between 10% to 50% of the total number of traces, trying all slices by range of 10%. The value of 50% corresponds to the original “full” training dataset composed of the first half of the trail of traces. We choose to use the stationary distance (*cf.* Section 4.1) as a direct measure of the similarity between the MMC learnt on the reduced training set and the one learnt on the full training set. More precisely, a small stationary distance between the reduced training MMC and the full training MMC means that the two models are quite similar and thus that there was enough information contained in the considered training set to build a “representative MMC”. From Figure 4, we can observe that the more traces are used to build the training MMC, the smaller the stationary distance becomes between the reduced training MMC and the full training MMC. In particular, it seems that we need at least 30% of the total of the mobility traces to build a compact and representative MMC model.

In the following, we first focus on the Geolife dataset in order to analyze and understand the behavior of the de-anonymizers and distances.

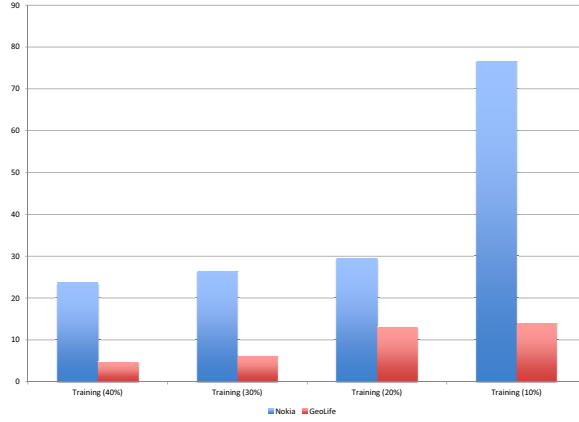


Figure 4: Stationary distance between the “reduced” training set and the “full” training set for the Geolife and the Nokia datasets, when the size of the training set varies between 10% to 50% of the total number of traces.

6.2. Fine-tuning clustering algorithms and de-anonymizers

The states of a MMC are extracted by running a clustering algorithm on the mobility traces of an individual. Therefore, the MMC generated (and by extension the success of the de-anonymization attack) is highly dependent on the clustering algorithm used and the accuracy of this algorithm, which may itself vary depending on the values of its parameters. The first step of our analysis consists in determining the parameters that leads to the best accuracy for the DJ-clustering algorithm.

Depending on the chosen values for the parameters, not all users of the original Geolife dataset will lead to the generation of a well-formed MMC. For instance, when the parameters of the clustering algorithm are too “conservative”, some users will not have enough mobility traces to identify frequent POIs, which results in their MMC being composed of only one state. On the contrary, choosing parameters that are too “relaxed” leads to the identification of a high number of POIs thus conducting to a MMC with too many states, which is detrimental to the success of the de-anonymization attack. Thus, the main objective of the tuning phase is to find the good set of parameters for the clustering algorithm that maximize the number of users in the training set whose MMCs does not consist in only one state while keeping the average number of POIs identified per user in an acceptable range.

First, we vary the three parameters of the clustering algorithm (*MinPts*, *r* and the minimal number of days) and count the number of MMC generated that have more than one state. We found that these parameters are themselves highly correlated with the duration of the collection period and the sampling rate used. Table 5 summarizes the values of the clustering parameters used in the following experiments obtained after this validation process.

Data set	<i>MinPts</i>	<i>r</i> (km)	Minimal nb of days
Arum	40	0.05	30
Geolife	20	0.5	10
Nokia	10	0.05	10
SFC	20	0.05	10
Borlange	2	0.05	10

Table 5: Validated clustering parameters.

Once the parameters of the clustering algorithm have been tuned, we have assessed the efficiency of the simple and weighted voting de-anonymizers. In particular for the simple voting method, we first studied the influence of

the number of candidates proposed by the minimal distance de-anonymizer used on each distance metric (stationary, matching, proximity and density-based). Each instance of this de-anonymizer proposes the n candidates that have the smallest distances in a row sorted in increasing order. The simple vote de-anonymizer is applied to this list of candidates in order to output a single prediction. Figure 5 illustrates the success rate of this attack as a function of the number of candidates (more precisely as the percentage of true positives obtained over the total number of individuals) and the number of users that have well-formed MMCs (*i.e.*, MMCs that have more than one state). This experiment was conducted on the Geolife dataset with the sampling rate varying in the range $\{10s, 30s, 60s, 120s\}$.

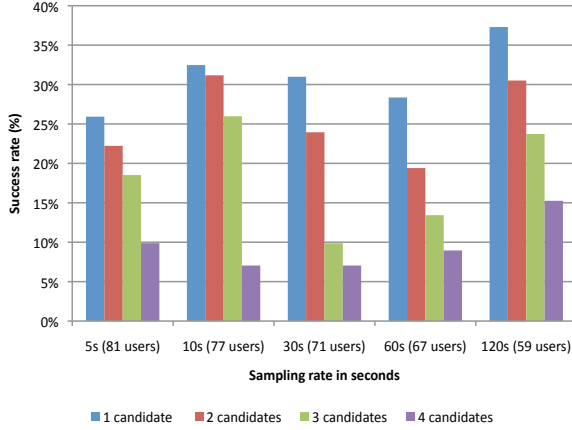


Figure 5: Success rate with the simple vote function.

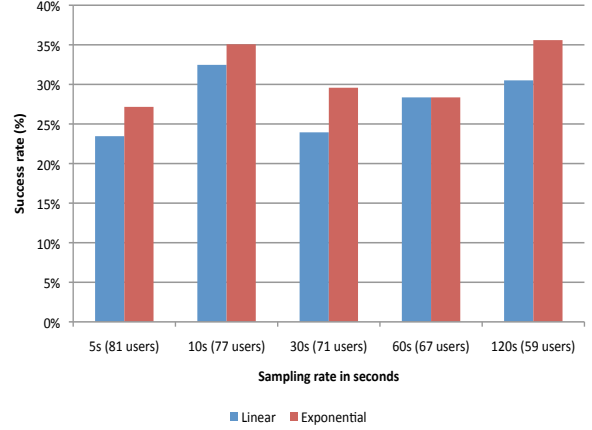


Figure 6: Success rate of the linear/exponential weighted votes.

From Figure 5, we can observe that the number of users considered that have a well-formed MMC decreases as the sampling rate increases, because a low sampling rate results in fewer mobility traces to build MMCs. We can also observe that the success rate of the attack decreases as the number of candidates increases, which is not surprising as a higher number of candidates renders the task of the de-anonymizer more complex than when they are few candidates. For instance, the success rate of the attack when 4 candidates are generated by each distance metric is never more than 15%. Therefore, we can conclude that for the simple vote method considering only one candidate per distance is sufficient.

However, in some situations it is helpful to consider more than one candidate per distance metric but their weight should be set to be different values, which is the main idea behind the weighted vote de-anonymizer. In our experiments, we have compared two different ways to weight candidates, one based on linear weights and the other on exponential ones to determine the most effective one (Figure 6). In a nutshell, the linear method assigns weights in a decreasing linear form and the exponential method assigns weights in a decreasing exponential form starting at 2^n , n being the number of candidates. Table 6 illustrates the assigned weights for these two methods, while Figure 6 compares the success rate of the weighted vote de-anonymizer using both weighting systems for different sampling rates. From these experiments, we can observe that the exponential method seems to be more efficient as its success rate is about 5% better than with the linear method. Moreover, this de-anonymizer seems to be robust to data sampling with different rates.

Candidate	Linear weight	Exponential weight
1	n	2^n
2	$n - 1$	2^{n-1}
...
n	1	2

Table 6: The linear and exponential weighting schemes.

6.3. Measuring the efficiency of de-anonymizers

To measure the success rate of the proposed de-anonymizers, we have sampled the Geolife dataset at different rates and observed the influence of the sampling on the success rates of the de-anonymizers. Figure 7 shows that the success rate of the attack with the minimal stationary distance and the minimal proximity distance varies from 20% to 40%, but that the best performing predictor is stat-prox with results ranging from 35% to 40%. By studying how successfully identified users tend to be the same across the different experiments, we have observed that a significant fraction of them tends to be re-identified with success independently of the sampling conditions. We believe that for these particular users their mobility behaviors is so regular that they are highly resilient to downsampling.

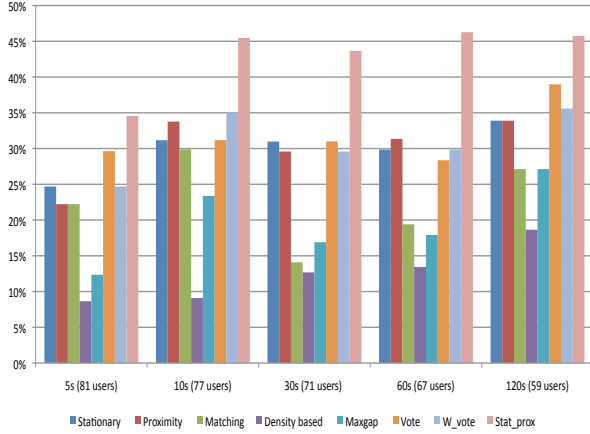


Figure 7: Success rate of the different de-anonymizers on the Geolife dataset.

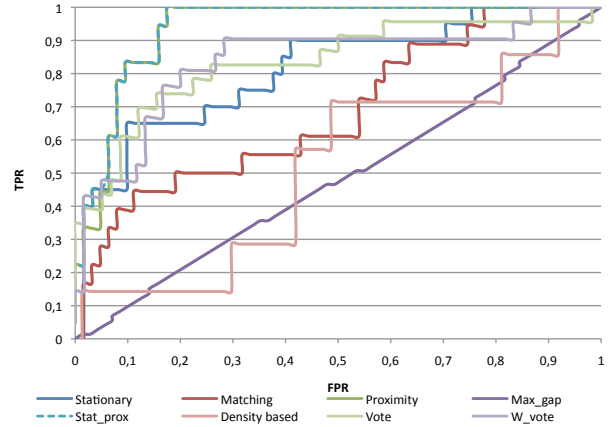


Figure 8: ROC curve for the different de-anonymizers on the Geolife dataset.

At this point of the experiments, it seems important to be able to compare precisely the de-anonymizers. Indeed, the success rate of a de-anonymization attack is not the only aspect that should be considered. For instance, for an adversary a possible strategy is to focus on weak individuals that offer a high probability of success for the attack rather than being able to de-anonymize the entire dataset. Measuring the probability of success of the inference attack for a given individual is similar to have some kind of confidence measure for a given de-anonymization candidate. Deriving this confidence measure is quite intuitive for our de-anonymizers. Indeed, for the minimal distance ones, the smaller is the distance, the higher the confidence.

In order to compare the performance of the de-anonymizers, we rely on the notion of *Receiver Operating Characteristic* (ROC) curve [33]. In a nutshell, a ROC curve is a graphical plot representing the sensitivity (*i.e.*, as measured by the true positives rate versus false positives rate) for a classifier. In our case, the ROC curve is built based on the confidence of the de-anonymizer. More precisely, for each candidate of the testing set, the distance between this candidate and each of the MMCs of the training set is computed. Afterwards, these distances are sorted by increasing order with the idea that the smaller the distance (or the higher the votes) the more confidence one can have in the result. These results sorted by confidence are then used to build the ROC curve in order to have the true positives at the beginning. The intuition behind this curve is that between two de-anonymizers achieving the same success rate, one should favor the one displaying the highest confidence. In Figure 8, the ROC curve shows the true positives rate (TPR) versus the false positives rate (FPR) for the best performing de-anonymizers, with the candidates sorted by ascending distance. This ROC curve further confirms that the stat-prox de-anonymizer is the best alternative among the de-anonymizers we designed.

Our approach performs fairly well for the Geolife dataset as the achieved success rate is between 35% and 45% for the stat-prox de-anonymizer (Figure 7). In order to further validate the approach, we applied it on the Nokia dataset. This dataset has 195 users, among which we can generate a “valid” MMC composed of more than one POI for 157 users using the parameters described previously. As shown on Figure 9, the success rate varies between 35% and 42%, with the best score obtained again by the stat-prox de-anonymizer.

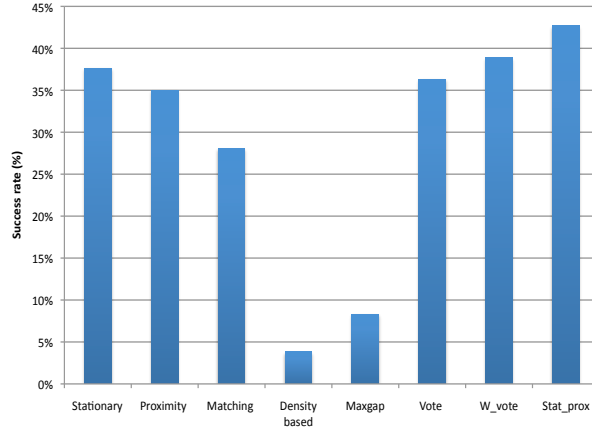


Figure 9: Success rate of the de-anonymizers on the Nokia dataset.

6.4. Fair comparison with prior work

In this section, we have presented various experiments on de-anonymization attacks that lead to the definition of an heterogeneous de-anonymizer called stat-prox, which obtains a success rate between 42% and 45% on different datasets. While at first glance, this performance may seem to be poorer than the one achieved by the predictors of Ma *et al.* [16], which goes up to 60% to 90%, and the predictors of Freudiger and co-authors [17] that have a re-identification rate of 70%, we believe that these results are not directly comparable because we clearly differentiate between the training set and the testing set, while these authors perform the learning and the testing on the same dataset, thus inducing a strong experimental bias.

Indeed, our mobility models are built out of the training set, which is disjoint from the test set, whereas one of the adversary model of Ma *et al.* directly extracts mobility traces forming the test set from the training set. Moreover in our case, the training data is temporally separated from the test data (*i.e.*, the training and the test have been recorded at different non-overlapping periods of time) because the whole dataset has been split into two temporally disjoint parts, whereas the second adversary model of Ma *et al.* picks the information it uses to de-anonymize within the same period as the test data is recorded. Therefore, our approach is quite different from them as our attack consists first in collecting mobility data from an individual, before trying to identify this individual in a so-called anonymized dataset, while their attack aims at gathering location data at the same time at which the de-anonymization attack occurs. In addition, one important parameter of their attack is the number of timestamped location data collected, which can be compared to the number of states we have in our mobility model. On average and depending on the dataset considered, we have between 4 and 8 states per MMC, which correspond to a compact representation of the mobility behavior of an individual. When restricted to such limited of information in terms of the number of timestamped location data, the attacks proposed by Ma *et al.* do not perform well, with a de-anonymization rate between 10% to 40%. Similarly in the work of Freudiger and co-authors, mobility samples are taken from home/work, POIs or frequent visited places without any distinction between training or testing sets, which again introduces an evaluation bias. For instance, when 90% of the samples are taken either from home or work, the success rate of the re-identification is around 75% while a uniform random sampling leads to a success rate for the re-identification of 10%.

For comparison purpose, we conducted the de-anonymization attack without separating the training and testing sets. The results obtained by related work and for the stat-prox de-anonymizer in this setting using different datasets are summarized in Figure 10. These experiments are, as expected, so biased that they lead to a success rate close to 100% for all the datasets. Once again, we do not pretend that our de-anonymization attack would achieve a success rate of nearly 100% and beat all the previous methods if tested in the same conditions (for instance in some settings the test set was only a subset of the training set as it was sampled from it). We are merely pointing out that for fairness issues, it is important to compare de-anonymizers using the same setting and that in order to reduce the experimental bias, the training and testing set should be clearly separated (which is not the case in almost all the previous works).

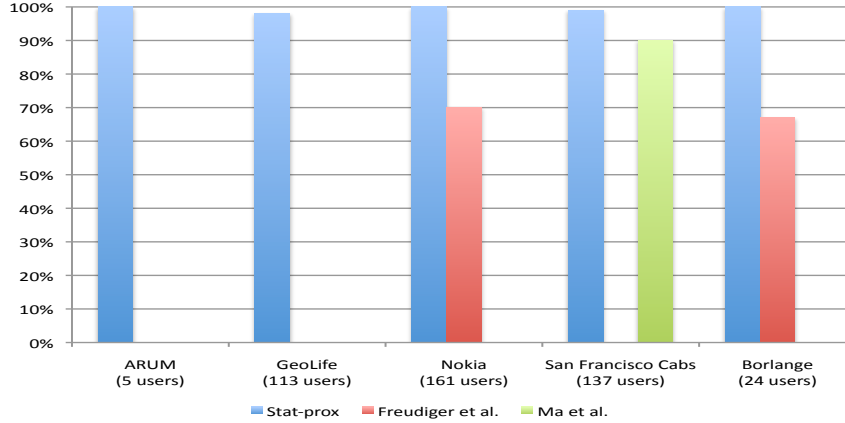


Figure 10: Success rate of the stat-prox de-anonymizer when the training and testing sets are the same.

7. Conclusion

In recent years, several privacy breaches related to location data have reached the headlines. For instance, the German deputy Malte Spitz sued Deutsche Telekom to obtain the last six months of location data generated from his phone [34]. Then, he published this data in the form of an interactive map showing that the combination of location data with contextual information could lead to a serious privacy breach. Another example of privacy buzz was the article about telephone constructors [35] published in The Wall Street Journal revealing that Apple and Google collect on a large scale location data using unique identifiers in order to develop novel location-based services.

The classical argument used by data collectors is that by itself location data is anonymous and thus can be collected from users without violating their privacy. Unfortunately, as shown by our work in this paper this argument is fallacious. More precisely, we have demonstrated that geolocated datasets gathering the movements of individuals are particularly vulnerable to a form of inference attack called the de-anonymization attack. More precisely, we have shown that the de-anonymization attack can re-identify with a high success rate the individuals whose movements are contained in an anonymous dataset provided that the adversary can use as background information some mobility traces of the same individuals that he has been able to observe during the training phase. From these traces, the adversary can build a MMC that models in a compact and precise way the mobility behavior of an individual. We designed novel distances quantifying the similarity between two MMCs and we described how these metrics can be combined to build de-anonymizers. The de-anonymization attack is very accurate with a success rate of up to 45% on large-scale real datasets and this even if the mobility traces are sanitized by downsampling them (*e.g.*, every 2 minutes instead of every 10 seconds).

In summary, the mobility behavior of an individual is far from being random [9] and tends to be unique thus acting as a signature of an individual [27]. For instance, even the pair home/work could act as a quasi-identifier [7]. In addition to location data, the knowledge of the social network, as shown by the works of Srivatsa and Hicks and Sharad and Danezis, can be used as side information to help in performing the de-anonymization. However, it might be possible to mitigate this risk of re-identification by sanitizing the social graph before releasing it [36].

In the future, we are planning to extend the current work by following several avenues of research. For instance, one of our research objective is to discover among different clustering algorithms, the one that best fits our needs while being also robust with respect to small changes in the inputs (*e.g.*, small spatial and temporal perturbation). In a different direction, we will also explore how more complex geo-sanitization mechanisms, such as spatial cloaking techniques [37] or mix zones [38], can help to reduce the success rate of the de-anonymization attack.

References

- [1] M. Piorkowski, N. Sarafijanovic-Djukic, M. Grossglauser, CRAWDAD data set epfl/mobility (v. 2009-02-24), Downloaded from <http://crawdad.cs.dartmouth.edu/epfl/mobility> (February 2009).

- [2] N. Kiukkonen, Data collection campaign, Tech. rep., Nokia Research Center, Lausanne, Switzerland (December 2009).
- [3] Y. Zheng, X. Xie, W.-Y. Ma, Geolife: A collaborative social networking service among user, location and trajectory, in: *IEEE Data Engineering Bulletin*, Vol. 33, Beijing, China, 2010, pp. 32–40.
- [4] M. Killijian, M. Roy, G. Trédan., Beyond San Francisco cabs: building a *-lity mining dataset, in: *Proceedings of Workshop on the Analysis of Mobile Phone Networks (NetMob)*, Cambridge, MA, USA, 2010, pp. 75–78.
- [5] J. Krumm, Inference attacks on location tracks, in: *Pervasive Computing*, Toronto, Canada, 2007, pp. 127–143.
- [6] Y. De Mulder, G. Danezis, L. Batina, B. Preneel, Identification via location-profiling in GSM networks, in: *Proceedings of the ACM Workshop on Privacy in the Electronic Society*, Alexandria, VA, USA, 2008, pp. 23–32.
- [7] S. Gams, M.-O. Killijian, M. Núñez del Prado Cortez, Show me how you move and I will tell you who you are, *Transactions on Data Privacy* 2 (4) (2011) 103–126.
- [8] H. Zang, J. C. Bolot, Anonymization of location data does not work: A large-scale measurement study, in: *Proceedings of ACM Mobicom*, Las Vegas, NV, USA, 2011, pp. 145–156.
- [9] M. C. Gonzalez, C. A. Hidalgo, Albert-Laszlo, Understanding individual human mobility patterns, in: *Nature*, New Orleans, LA, USA, 2008, pp. 779–782.
- [10] C. Song, Z. Qu, N. Blumm, A.-L. Barabási, Limits of predictability in human mobility, in: *Science*, Vol. 327, New Orleans, LA, USA, 2010, pp. 1018–1021.
- [11] L. Jedrzejczyk, B. Price, A. Bandara, B. Nuseibeh, I know what you did last summer: risks of location data leakage in mobile and social computing, Tech. rep., Department of Computing Faculty of Mathematics, Computing and Technology The Open University, Milton Keynes, UK (November 2008).
- [12] C. Bettini, X. S. Wang, S. J. Jodia, Protecting privacy against location-based personal identification, *Privacy and Security Support for Distributed Applications* 3674 (8) (2005) 185–199.
- [13] A. Narayanan, V. Shmatikov, Robust de-anonymization of large sparse datasets, in: *Proceedings of the IEEE Symposium on Security and Privacy*, Washington, DC, USA, 2008, pp. 111–125.
- [14] P. Tan, M. Steinbach, V. Kumar, *Introduction to data mining*, Pearson International Edition, Pearson Addison Wesley, Boston, 2006.
- [15] N. Eagle, A. Sandy Pentland, Reality mining: sensing complex social systems, in: *Personal and Ubiquitous Computing*, Vol. 10, London, UK, 2006, pp. 255–268.
- [16] C. Y. Ma, D. K. Yau, N. K. Yip, N. S. Rao, Privacy vulnerability of published anonymous mobility traces, in: *Proceedings of the 16th annual international conference on Mobile computing and networking*, New York, NY, USA, 2010, pp. 185–196.
- [17] J. Freudiger, R. Shokri, J.-P. Hubaux, Evaluating the privacy risk of location-based services, in: *Proceedings of the 15th international conference on Financial Cryptography and Data Security*, Berlin, Germany, 2012, pp. 31–46.
- [18] X. Xiao, Y. Zheng, Q. Luo, X. Xie, Finding similar users using category-based location history, in: *Proceedings of the International Conference on Advances in Geographic Information Systems*, New York, NY, USA, 2010, pp. 442–445.
- [19] Y. Zheng, Q. Li, Y. Chen, X. Xie, W.-Y. Ma., Understanding mobility based on GPS data., in: *In Proceedings of ACM conference on Ubiquitous Computing*, Seoul, Korea, 2008, pp. 312–321.
- [20] K. Järvelin, J. Kekäläinen, Cumulated gain-based evaluation of ir techniques, *ACM Transactions on Information Systems* 20 (4) (2002) 422–446.
URL <http://doi.acm.org/10.1145/582415.582418>
- [21] R. Shokri, G. Theodorakopoulos, J.-Y. Le Boudec, J.-P. Hubaux, Quantifying Location Privacy, in: *Proceedings of IEEE Symposium on Security and Privacy (S&P)*, 2011.
- [22] S. J. Russell, P. Norvig, J. F. Candy, J. M. Malik, D. D. Edwards, *Artificial intelligence: a modern approach*, Prentice-Hall, Inc., Upper Saddle River, NJ, USA, 1996.
- [23] M. Srivatsa, M. Hicks, De-anonymizing mobility traces: using social network as a side-channel, in: *Proceedings of the ACM conference on Computer and communications security*, New York, NY, USA, 2012, pp. 628–637.
- [24] P. Crucitti, V. Latora, S. Porta, Centrality measures in spatial networks of urban streets, *Physical Review E* 73 (3) (2006) 036125.
- [25] K. Sharad, G. Danezis, De-anonymizing d4d datasets, in: *Workshop on Hot Topics in Privacy Enhancing Technologies*, Bloomington, Indiana, USA, 2013.
- [26] J. Leskovec, J. Kleinberg, C. Faloutsos, Graph evolution: Densification and shrinking diameters, *Trans. Knowl. Discov. Data* 1 (1).
- [27] Y.-A. de Montjoye, C. A. Hidalgo, M. Verleysen, V. D. Blondel, Unique in the Crowd: The privacy bounds of human mobility, *Scientific Reports* 3 (3) (2013) 1 – 5.
- [28] J. Unnikrishnan, F. M. Naini, De-anonymizing private data by matching statistics, in: *Conference on Communication, Control, and Computing*, Monticello, Illinois, USA, 2013.
- [29] M. Gutman, Asymptotically optimal classification for multiple tests with empirically observed statistics, *Transactions on Information Theory* 35 (2) (1989) 401–408.
- [30] D. Ashbrook, T. Starner, Learning significant locations and predicting user movement with GPS, in: *Proceedings of the IEEE International Symposium on Wearable Computers*, Sardina, Italy, 2003, pp. 275–286.
- [31] Z. Yan, D. Chakraborty, C. Parent, S. Spaccapietra, K. Aberer, SeMiTri: a framework for semantic annotation of heterogeneous trajectories, *Proceedings of the International Conference on Extending Database Technology* 5 (1) (2011) 259–270.
- [32] J. Munkres, Algorithms for the Assignment and Transportation Problems, *Journal of the Society for Industrial and Applied Mathematics* 5 (1) (1957) 32–38.
- [33] T. Fawcett, An introduction to ROC analysis, *Pattern Recognition Letters* 27 (8) (2006) 861–874.
- [34] Z. online, Tell-all telephone, <http://www.zeit.de/datenschutz/malte-spitz-data-retention/>.
- [35] J. V.-D. Julia Angwin, Apple, google collect user data, <http://online.wsj.com/article/SB10001424052748703983704576277101723453610.html>.
- [36] G. Gulys, S. Imre, Hiding information in social networks from de-anonymization attacks by using identity separation, in: *Proceedings of the Joint IFIP TC6 and TC11 Conference on Communications and Multimedia Security*, Magdeburg, Germany, 2013, pp. 173–184.

- [37] M. Gruteser, D. Grunwald, Anonymous usage of location-based services through spatial and temporal cloaking, in: Proceedings. of the international conference on Mobile systems, applications and services, New York, NY, USA, 2003, pp. 31–42.
- [38] A. R. Beresford, F. Stajano, Location privacy in pervasive computing, *IEEE Pervasive Computing* 2 (1) (2003) 46–55.